

# Automated Verification of Security Chains in Software-Defined Networks with Synaptic

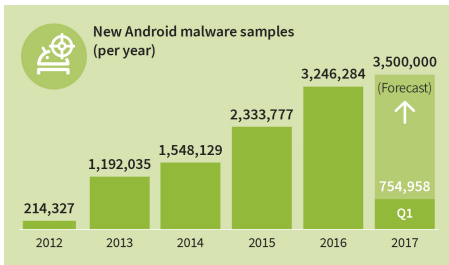
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- 1 Problem Statement
- 2 Related Work
- 3 Automated Verification of Security Chains
- 4 Prototype and Performance Evaluation
- 5 Conclusions

- Growing number of malicious applications on Android devices
  - 3.25 million malicious applications in 2016 (Gdata Security Labs)

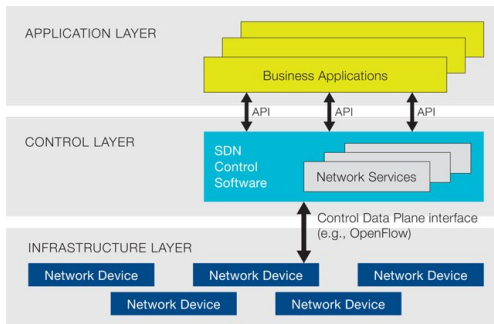


- Limitations of market preventive security methods
- On-device security constrained in terms of resources
- Partial outsourcing through cloud-based security chains
- Can automated verification methods help ensure their quality?

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- Virtualization and outsourcing of network security functions in cloud infrastructures
- Target: higher flexibility and better manageability
- Uniform middlebox modelling proposed by Joseph and Stoica
- Firewall middlebox specification with *FlowGuard*
- Automating large-scale deployments remains difficult
- Lack of uniform interfaces for these devices
- Formal verification could help validating their correctness

# Network programmability as a management enabler



SDN architecture (source: [www.opennetworking.org](http://www.opennetworking.org))

- Easier automation of middlebox deployment and chaining
- Low level interface needs abstraction and formal methods

## SMT solving

- Representation of systems and properties as logical formulas
- SMTlib: a standard input language for SMT solvers
- Verification based on checking satisfiability/validity

## Model checking

- Representation of systems based on finite state automata
- Properties expressed in temporal logics: LTL (Linear Temporal Logic) or CTL (Computation Tree Logic)

## Chaining security functions for protecting end users

- Automated building in cloud infrastructures using SDN
- Dynamic and optimized placement of security functions based on available resources

## Efforts for formally verifying such security compositions

- Validation of distributed firewall policies: *Al-Shaer et al.*
  - Example: do not simultaneously authorize and block port 5000
- Validation of traffic properties: *Vericon (Ball et al.)*
  - Example: check that a policy blocks incoming connections



## Pyretic

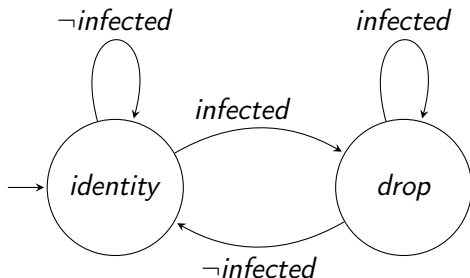
- Python-embedded domain-specific programming language for writing SDN control programs
- Network data-plane behavior: policy functions
- compiles policies to OpenFlow-based switches

## Pyretic rules and compositions

- 4 atomic rules: identity, drop, match and modify
- Sequential composition operator: `»`
- Parallel composition operator: `+`

## Examples of simple policy chains

- Parallel composition: `match(port=5000) + match(port=4000)`
- Sequential composition: `match(port=5000) » drop`

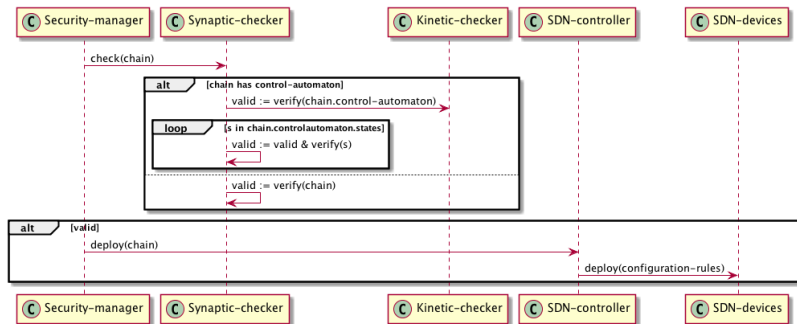


- **Kinetic**: represent control dynamics as a Finite State Machine
- Validation by model checking with NuSMV
- Example property: **AG** (*infected* → **AX** *drop*)
- Objective: extend formal verification to the data plane

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## Verification of both control and data planes of security chains

- Extension of the Kinetic language to the data plane
- Verification using SMT solving and model checking
  - Chains translated to both types of tools
  - Addition of properties to constrain the model



## Modelling and translation

- Chains and constraints expressed with SMTlib
- Representation based on logical formulas
- Prototype supporting the CVC4 and veriT SMT solvers

## Advantages of SMT solving for our strategy

- Pyretic expressions correspond to logical formulas
- Easy to add constraints expressed as equalities
- Example: `port==5000`  $\rightarrow$  flow accepted

# Compilation of a chain for SMTlib

- Atomic rules represented using Boolean values and equalities
- These rules represent the basic building blocks
- Composition operators encoded using Boolean connectives
- Used to build up complex security chains
- Traffic invariants expressed as logical conditions
- Success: the chain validates the constraints

## Compilation of a simple chain

- `match(port=5000) + match(port=4000)`  
→ `(OR (= port 5000) (= port 4000))`
- Production of a packet condition

## Compilation of a condition

- `ip=192.168.0.15` → `(= ip 1)`
- replacement of network elements by integer values

## Modelling and translation

- Represent the network as a finite state automaton
- Express the invariants in temporal logic (CTL)
- Success: the automaton satisfies the invariants

## Advantages of model checking for our strategy

- Already used for the control plane (Kinetic)
- Possibility of extending this model to the data plane by using nuXmv instead of NuSMV



- Express the network as a finite state automaton
- Variables represent packet fields
  
- Strictly sequential policies correspond to automaton transitions
- Join points give rise to states of the automaton
  
- Traffic invariants expressed as CTL properties
- Authorized traffic accepted by the automaton
- Forbidden packets are blocked

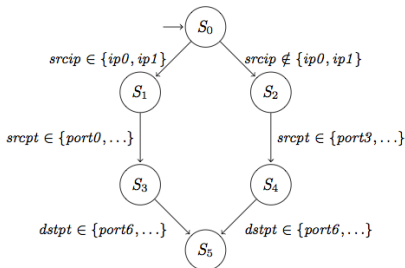
# Example of verification by model checking

```
F1 = match(srcip=IP("198.122.37.15")) +  
     match(srcip=IP("253.182.3.14"))  
  
F2 = match(srcport=1000) + match(srcport=2000) +  
     match(srcport=3000)  
  
F3 = match(srcport=4000) + match(srcport=5000) +  
     match(srcport=6000)  
  
F4 = match(dstport=7000) + match(dstport=8000) +  
     match(dstport=9000)  
  
chain = ((F1 >> F2) + (~F1 >> F3)) >> F4
```

Pyretic firewall chain

```
AG(((srcip=ip0 | srcip=ip1) &  
    (srcpt=port0 | srcpt=port1 | srcpt=port2) &  
    (dstpt=port6 | dstpt=port7 | dstpt=port8))  
-> EF state=S5)  
  
AG(((!(srcip=ip0 | srcip=ip1)) &  
    (srcpt=port3 | srcpt=port4 | srcpt=port5) &  
    (dstpt=port6 | dstpt=port7 | dstpt=port8))  
-> EF state=S5)
```

Chain specifications in CTL



Data plane FSM

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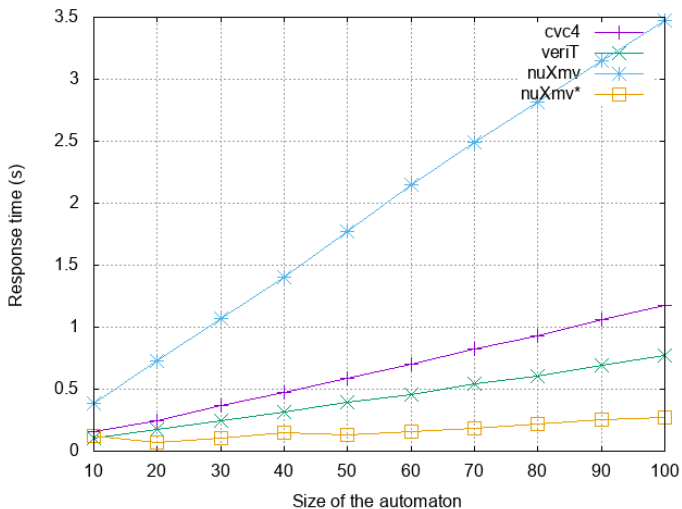
# Our prototype: the Synaptic checker

- Checker implemented in Python
- Extension of Pyretic and Kinetic languages
- Implementation of a module dedicated to traffic properties
  
- Integration of formal verification methods
- Use nuXmv or SMT solvers as backend verifiers

## Evaluation based on a set of synthetic security chains

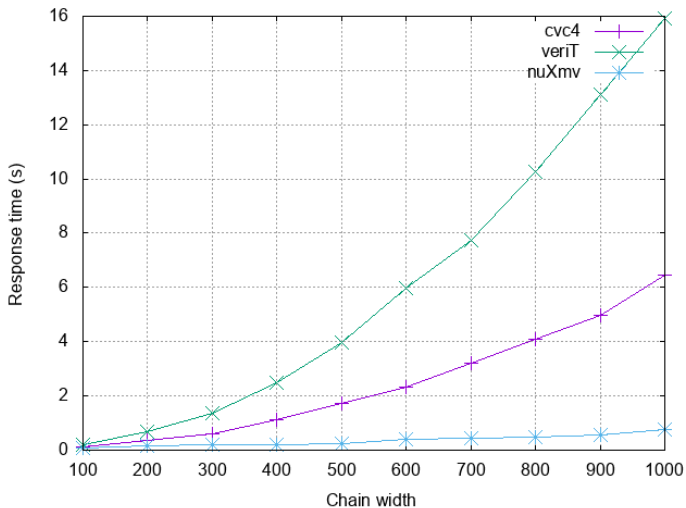
- Impact of several chain-related parameters
  - Size of the control automaton
  - Length and width of a security chain
  - Number of properties to be verified
- Comparison of verification performances
  - Response time
  - Memory space consumption

# Impact of the size of the control plane automaton



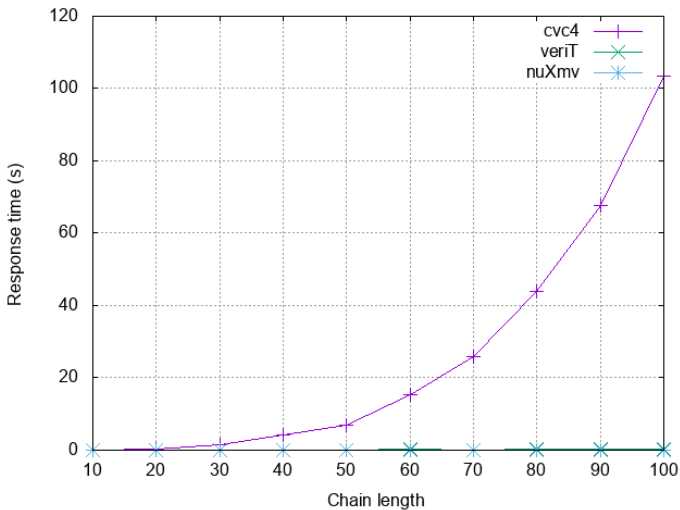
Response time vs. size of the control automaton

# Impact of the width of a chain



Response time vs. width of a chain

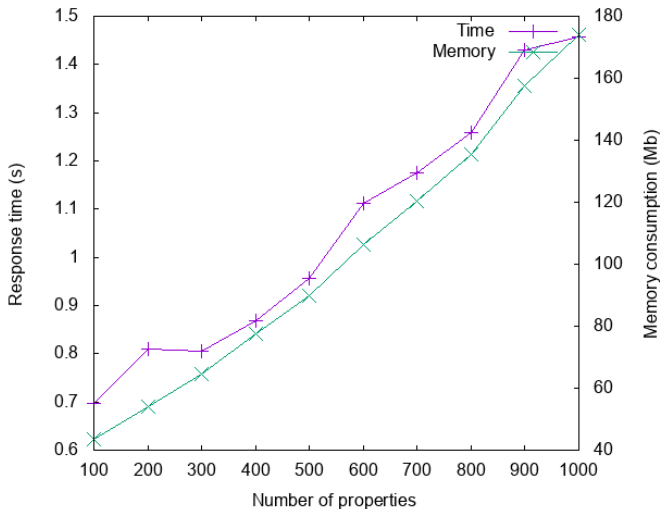
# Impact of the length of a chain



Response time vs. length of a chain



# Impact of the number of properties to verify



Resource consumption vs. number of properties (nuXmv)

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## Automated techniques for verifying SDN security chains

- Checking both the control and the data planes
- Translation of security chains for automatic verification
- Backends based on SMT solving or model checking

## Prototyping and performance evaluation

- Development of the Synaptic checker
- nuXmv shows best overall performance in our experiments

## Perspectives

- Extensions for more complex and advanced rules
- Further optimization of the encodings into SMTlib and nuXmv
- Integration into an automated security framework